

Linear Time Admission Control for Elastic Scheduling

Work Already Published

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Overview

- Elastic scheduling is a model for compressing task utilizations in an overloaded system
- The original algorithm* performs compression in time quadratic in the number of tasks
- In this work, we present an implementation that, with $O(n \log n)$ initialization, performs compression for online task admission in linear time

* Giorgio C. Buttazzo, Giuseppe Lipari, and Luca Abeni. “Elastic task model for adaptive rate control.” RTSS 1998.

Review of the Original Algorithm

Each task τ_i characterized by:

- U_i^{max} : Initial, uncompressed utilization
- E_i : Elasticity, flexibility to vary utilization
- U_i^{min} : Constraint on minimum utilization

System global variables:

- E_{SUM} : Total elasticity (sum of E_i)
- U_{SUM} : Initial total utilization (sum of U_i^{max})
- U_d : Desired total utilization

During overload, total required compression is

$$U_{SUM} - U_d$$

Algorithm

1. Compress each task's utilization proportionally to its elasticity:

$$U_i = U_i^{max} - (U_{SUM} - U_d) \frac{E_i}{E_{SUM}}$$

2. While any task τ_i has $U_i < U_i^{min}$:
 - a. Set $U_i = U_i^{min}$
 - b. Repeat from 1

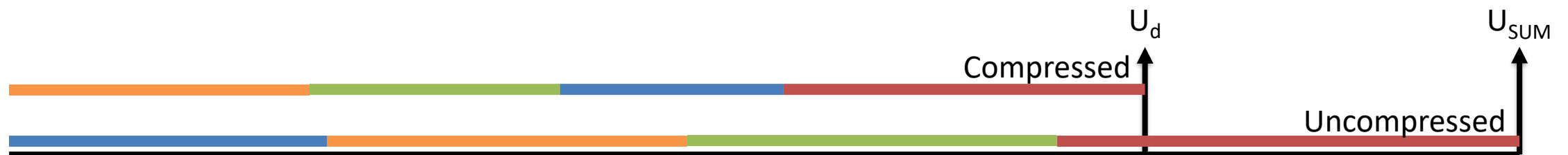
Improved Compression Algorithm

Initialization: Sort tasks in a list according to $\phi_i = \frac{U_i^{max} - U_i^{min}}{E_i}$

Task Admission:

1. Insert new task τ_j according to ϕ_j
2. For each task τ_i :

- a. Compress according to $U_i = U_i^{max} - (U_{SUM} - U_d) \frac{E_i}{E_{SUM}}$
- b. Is $U_i < U_i^{min}$?
- c. If so, set $U_i = U_i^{min}$
- d. If not, no remaining task will have $U_i < U_i^{min}$, and so compress normally



Conclusions

- We have presented a new algorithm for elastic compression, which improves on the prior $O(n^2)$ algorithm:
- It requires one-time $O(n \log n)$ initialization to sort tasks
- Enables new task admission and compression in $O(n)$ time
- If a task leaves, it is removed from the (still sorted) list, allowing decompression in $O(n)$ time
- Questions? msudvarg@wustl.edu
- Sudvarg, M., Gill, C. & Baruah, S. Linear-time admission control for elastic scheduling. *Real-Time Systems* **57**, 485–490 (2021)